Weighted spanning tree enumerators of complete colorful complexes

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CombinaTexas University of Houston – Downtown April 21, 2013

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Spanning trees of K_n

Theorem (Cayley) K_n has n^{n-2} spanning trees.

 $T \subseteq E(G)$ is a **spanning tree** of G when:

- 0. spanning: T contains all vertices;
- 1. connected $(\tilde{H}_0(T) = 0)$
- 2. no cycles $(\tilde{H}_1(T) = 0)$
- 3. correct count: |T| = n 1

If 0. holds, then any two of 1., 2., 3. together imply the third condition.

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Theorem (Cayley-Prüfer)

$$\sum_{T\in ST(K_n)} \text{wt } T = (x_1 \cdots x_n)(x_1 + \cdots + x_n)^{n-2},$$

where wt $T = \prod_{e \in T}$ wt $e = \prod_{e \in T} (\prod_{v \in e} x_v)$.

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Example (K_4)

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Example (K_4)

► 4 trees like:
$$T = 2$$
 ↓ ↓ wt $T = (x_1 x_2 x_3 x_4) x_2^2$

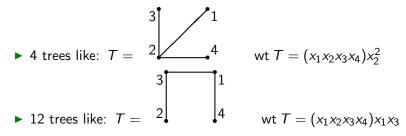
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Example (K_4)



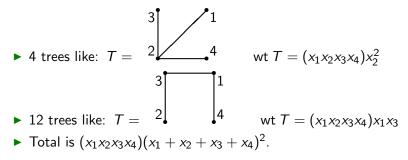
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Example (K_4)



Counting trees Matrix-tree theorem

Complete bipartite graphs

Example $(K_{3,2})$

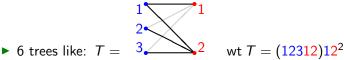
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Counting trees Matrix-tree theorem

Complete bipartite graphs

Example $(K_{3,2})$

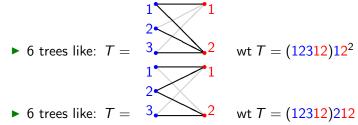


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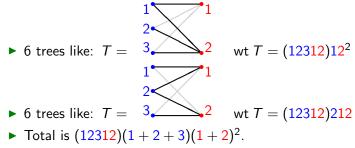


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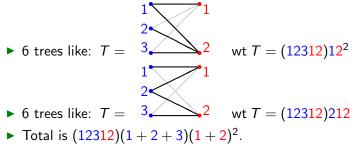


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Counting trees Matrix-tree theorem

Complete bipartite graphs

Example $(K_{3,2})$



Theorem

$$\sum_{T\in ST(K_{m,n})} \operatorname{wt} T = (x_1 \cdots x_m)(y_1 \cdots y_n)(x_1 + \cdots + x_m)^{n-1}(y_1 + \cdots + y_n)^{m-1}.$$

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Counting trees Matrix-tree theorem

Laplacian

Theorem (Kirchoff's Matrix-Tree)G has $|\det L_r(G)|$ spanning treesDefinition TheLaplacian matrix of graph G, denoted byL (G).

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Laplacian

Theorem (Kirchoff's Matrix-Tree) *G* has $|\det L_r(G)|$ spanning trees **Definition** The Laplacian matrix of graph *G*, denoted by *L* (*G*). Defn 1: L(G) = D(G) - A(G) $D(G) = \text{diag}(\deg v_1, \dots, \deg v_n)$ A(G) = adjacency matrix

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Laplacian

Theorem (Kirchoff's Matrix-Tree) G has $|\det L_r(G)|$ spanning trees **Definition** The Laplacian matrix of graph G, denoted by L (G). Defn 1: L(G) = D(G) - A(G) $D(G) = \text{diag}(\text{deg } v_1, \ldots, \text{deg } v_n)$ A(G) = adjacency matrixDefn 2: $L(G) = \partial(G)\partial(G)^T$ $\partial(G) =$ incidence matrix (boundary matrix)

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Laplacian

Theorem (Kirchoff's Matrix-Tree) G has $|\det L_r(G)|$ spanning trees **Definition** The reduced Laplacian matrix of graph G, denoted by $L_r(G)$. Defn 1: L(G) = D(G) - A(G) $D(G) = \text{diag}(\text{deg } v_1, \ldots, \text{deg } v_n)$ A(G) = adjacency matrixDefn 2: $L(G) = \partial(G)\partial(G)^T$ $\partial(G) =$ incidence matrix (boundary matrix) "Reduced": remove rows/columns corresponding to any one vertex

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Counting trees Matrix-tree theorem

Example $(K_{3,2})$

	1	1			11	12	21	22	31	<mark>3</mark> 2
	\sim			1		-1	0	0	0	0
3 2			a	$=$ $\frac{2}{3}$	0	0	-1	-1	0 -1	0
			0	3		0	0	0	-1	-1
				1	1	0	1	0	1	0
				2	0	1	0	1	0	1
	$\binom{2}{2}$	0	0	- 1	-1°)				
	0	2	0	-1	-1					
L =	0	0	2	-1	-1					
	-1	-1	-1	3	0					
	$ \begin{pmatrix} 2 \\ 0 \\ -1 \\ -1 \end{pmatrix} $	-1	-1	0	3	/				

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Counting trees Matrix-tree theorem

Example $(K_{3,2})$

	1	1			11						
	\sim			1	-1	-1	0	0	0	0	_
		$\searrow 2$	ລ	2	-1 0 0	0	-1	-1	0	0	
			0	$\partial = \frac{2}{3}$		0	0	0	-1	-1	
				1	1 0	0	1	0	1	0	
				2	0	1	0	1	0	1	
L =	$\begin{pmatrix} 2 \\ 0 \\ 0 \\ -1 \\ -1 \end{pmatrix}$	0 2 0 -1 -1	0 0 2 -1 -1		$ -1 \\ -1 \\ $	L_r	= (2 0 -1 -1	0 2 -1 -1	$-1 \\ -1 \\ 3 \\ 0$	$\begin{pmatrix} -1 \\ -1 \\ 0 \\ 3 \end{pmatrix}$

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Counting trees Matrix-tree theorem

Example $(K_{3,2})$

	1	1			11						
	\sim			1	-1	-1	0	0	0	0	-
		$\searrow 2$	∂	2	0	0	-1	-1	0	0	
	<i></i>		0	= 3	-1 0 0	0	0	0	-1	-1	
				1	1	0	1	0	1	0	
				2	1 0	1	0	1	0	1	
L =	$\begin{pmatrix} 2 \\ 0 \\ 0 \\ -1 \\ -1 \end{pmatrix}$	0 2 0 -1 -1	0 0 2 -1 -1			L_r		2 0 —1 —1	0 2 -1 -1	$-1 \\ -1 \\ 3 \\ 0$	$\begin{pmatrix} -1 \\ -1 \\ 0 \\ 3 \end{pmatrix}$

 $det(L_r) = 12$, the number of spanning trees of $K_{3,2}$.

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Counting trees Matrix-tree theorem

Weighted Matrix-Tree Theorem

$$\sum_{T\in ST(G)} \operatorname{wt} T = |\det \hat{L}_r(G)|,$$

where $\hat{L}_r(G)$ is reduced weighted Laplacian. Defn 1: $\hat{L}(G) = \hat{D}(G) - \hat{A}(G)$ $\hat{D}(G) = \operatorname{diag}(\operatorname{deg} v_1, \ldots, \operatorname{deg} v_n)$ $\hat{\deg v_i} = \sum_{v_i v_i \in E} x_i x_j$ $\hat{A}(G) = adjacency matrix$ (entry $x_i x_i$ for edge $v_i v_i$) Defn 2: $\hat{L}(G) = \partial(G)B(G)\partial(G)^T$ $\partial(G) =$ incidence matrix B(G) diagonal, indexed by edges, entry $\pm x_i x_i$ for edge $v_i v_i$

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Counting trees Matrix-tree theorem

Example $(K_{3,2})$



$$\hat{L}_r = \begin{pmatrix} 2(1+2) & 0 & -21 & -22 \\ 0 & 3(1+2) & -31 & -32 \\ -21 & -31 & 1(1+2+3) & 0 \\ -22 & -32 & 0 & 2(1+2+3) \end{pmatrix}$$
$$\det \hat{L}_r = (12312)(1+2+3)(1+2)^2$$

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Simplicial spanning trees of K_n^d [Kalai, '83]

Let K_n^d denote the complete *d*-dimensional simplicial complex on *n* vertices. $\Upsilon \subseteq K_n^d$ is a **simplicial spanning tree** of K_n^d when:

0.
$$\Upsilon_{(d-1)} = K_n^{d-1}$$
 ("spanning");

1.
$$\tilde{H}_{d-1}(\Upsilon; \mathbb{Z})$$
 is a finite group ("connected");

2.
$$\tilde{H}_d(\Upsilon; \mathbb{Z}) = 0$$
 ("acyclic");

3.
$$|\Upsilon| = \binom{n-1}{d}$$
 ("count").

- If 0. holds, then any two of 1., 2., 3. together imply the third condition.
- When d = 1, coincides with usual definition.

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C<mark>omplete skeleton</mark> Arbitrary complexes

 $= n^{\binom{n-2}{d}}$

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Counting simplicial spanning trees of K_n^d

Conjecture [Bolker '76]

$$\sum_{\Upsilon \in SST(K_n^d)}$$

Counting simplicial spanning trees of K_n^d

Theorem [Kalai '83]

$$\tau(K_n^d) = \sum_{\Upsilon \in SST(K_n^d)} |\tilde{H}_{d-1}(\Upsilon)|^2 = n^{\binom{n-2}{d}}$$

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Weighted simplicial spanning trees of K_n^d

As before,

wt
$$\Upsilon = \prod_{F \in \Upsilon}$$
 wt $F = \prod_{F \in \Upsilon} (\prod_{v \in F} x_v)$

Example

$$\begin{split} & \Upsilon = \{123, 124, 125, 134, 135, 245\} \\ & \text{wt} \ \Upsilon = x_1^5 x_2^4 x_3^3 x_4^3 x_5^3 \end{split}$$

Weighted simplicial spanning trees of K_n^d

As before,

wt
$$\Upsilon = \prod_{F \in \Upsilon}$$
 wt $F = \prod_{F \in \Upsilon} (\prod_{v \in F} x_v)$

Example

$$\begin{split} & \Upsilon = \{123, 124, 125, 134, 135, 245\} \\ & \text{wt} \ \Upsilon = x_1^5 x_2^4 x_3^3 x_4^3 x_5^3 \end{split}$$

Theorem (Kalai, '83)

$$\hat{\tau}(K_n^d) := \sum_{T \in SST(K_n^d)} |\tilde{H}_{d-1}(\Upsilon)|^2 (\text{wt } \Upsilon)$$

$$= (x_1 \cdots x_n)^{\binom{n-2}{d-1}} (x_1 + \cdots + x_n)^{\binom{n-2}{d}}$$

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Proof

Proof uses determinant of reduced Laplacian of K_n^d . "Reduced" now means pick one vertex, and then remove rows/columns corresponding to all (d-1)-dimensional faces containing that vertex.

$$\begin{split} L &= \partial \partial^T \\ \partial \colon \Delta_d \to \Delta_{d-1} \text{ boundary} \\ \partial^T \colon \Delta_{d-1} \to \Delta_d \text{ coboundary} \\ \text{Weighted version: Multiply column } F \text{ of } \partial \text{ by } x_F \end{split}$$

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Example n = 4, d = 2 (tetrahedron)

 $det L_r = 4$

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Simplicial spanning trees of arbitrary simplicial complexes

Let Δ be a *d*-dimensional simplicial complex. $\Upsilon \subseteq \Delta$ is a **simplicial spanning tree** of Δ when:

0.
$$\Upsilon_{(d-1)} = \Delta_{(d-1)}$$
 ("spanning");

1.
$$\tilde{H}_{d-1}(\Upsilon; \mathbb{Z})$$
 is a finite group ("connected");

2.
$$\tilde{H}_d(\Upsilon; \mathbb{Z}) = 0$$
 ("acyclic");

3.
$$f_d(\Upsilon) = f_d(\Delta) - \tilde{\beta}_d(\Delta) + \tilde{\beta}_{d-1}(\Delta)$$
 ("count").

- If 0. holds, then any two of 1., 2., 3. together imply the third condition.
- When d = 1, coincides with usual definition.

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Simplicial Matrix-Tree Theorem

Theorem (D.-Klivans-Martin, '09)

$$\hat{ au}(\Delta) = rac{| ilde{H}_{d-2}(\Delta;\mathbb{Z})|^2}{| ilde{H}_{d-2}(\Gamma;\mathbb{Z})|^2}\det\hat{L}_{\Gamma},$$

where

- ► $\Gamma \in SST(\Delta_{(d-1)})$
- $\partial_{\Gamma} = restriction of \partial_d$ to faces not in Γ
- reduced Laplacian $L_{\Gamma} = \partial_{\Gamma} \partial^{T}_{\Gamma}$
- Weighted version: Multiply column F of ∂ by x_F

Note: The $|\tilde{H}_{d-2}|$ terms are often trivial.

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Example: Octahedron

- Vertices 1, 2, 1, 2, 1, 2.
- ▶ Facets 111, 112, 121, 122, 211, 212, 221, 222,
- Γ = 11, 12, 11, 12, 22 spanning tree of 1-skeleton, so remove (rows and columns corresponding to) those edges from weighted Laplacian.
- det $\hat{L}_{\Gamma} = (121212)^3(1+2)(1+2)(1+2)$.

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Complete colorful complexes

Definition (Adin, '92)

The complete colorful complex $K_{n_1,...,n_r}$ is a simplicial complex with:

• vertex set $V_1 \dot{\cup} \dots \dot{\cup} V_r$ (V_i is set of vertices of color *i*);

$$|V_i| = n_i;$$

faces are all sets of vertices with no repeated colors.

Example

Octahedron is K_{222} .

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Unweighted enumeration

Theorem (Adin, '92)

The top-dimensional spanning trees of $K_{n_1,...,n_r}$ are "counted" by

$$au(K_{n_1,...,n_r}) = \prod_{i=1}^r n_i^{\prod_{j\neq i}(n_j-1)}.$$

Note: Adin also has a more general formula for dimension less than r-1.

Example

•
$$\tau(K_{222}) = 2^1 \times 2^1 \times 2^1$$

• $\tau(K_{235}) = 2^{2 \cdot 4} \times 3^{1 \cdot 4} \times 5^{1 \cdot 2}$
• $\tau(K_{mn}) = m^{n-1} \times n^{m-1}$

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Weighted enumeration

Theorem (Aalipour-D.)

The top-dimensional spanning trees of $K_{n_1,...,n_r}$ are "counted" by $\tau(K_{n_1,...,n_r}) =$

$$\prod_{i=1}^{r} (x_{i,1} + \cdots + x_{i,n_i})^{\prod_{j \neq i} (n_j - 1)} (x_{i,1} \cdots x_{i,n_i})^{(\prod_{j \neq i} n_j) - (\prod_{j \neq i} (n_j - 1))}.$$

Example

r

$$\begin{aligned} \hat{\tau}(\mathcal{K}_{235}) &= (x_1 + x_2)^{2 \cdot 4} (x_1 x_2)^{3 \cdot 5 - 2 \cdot 4} \\ &\times (y_1 + y_2 + y_3)^{1 \cdot 4} (y_1 y_2 y_3)^{2 \cdot 5 - 1 \cdot 4} \\ &\times (z_1 + \dots + z_5)^{1 \cdot 2} (z_1 \dots z_5)^{2 \cdot 3 - 1 \cdot 2} \end{aligned}$$

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Unweighted enumeration Weighted enumeration

Proof (via example $K_{3,2}$)

$$det \begin{pmatrix} 2(1+2) & 0 & -21 & -22 \\ 0 & 3(1+2) & -31 & -32 \\ -21 & -31 & 1(1+2+3) & 0 \\ -22 & -32 & 0 & 2(1+2+3) \end{pmatrix}$$
$$2312 det \begin{pmatrix} 1+2 & 0 & -1 & -2 \\ 0 & 1+2 & -1 & -2 \\ -2 & -3 & 1+2+3 & 0 \\ -2 & -3 & 0 & 1+2+3 \end{pmatrix}$$

By "identification of factors" (Martin-Reiner, '03), to show $(1+2)^2$ is a factor of the determinant, we just have to show that the nullspace of this matrix is at least 2, when we set 1+2=0.

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Finding null vectors

$$\begin{pmatrix} 1+2 & 0 & -1 & -2 \\ 0 & 1+2 & -1 & -2 \\ -2 & -3 & 1+2+3 & 0 \\ -2 & -3 & 0 & 1+2+3 \end{pmatrix}$$

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Finding null vectors

$$\begin{pmatrix} 1+2 & 0 & -1 & -2 \\ 0 & 1+2 & -1 & -2 \\ -2 & -3 & 1+2+3 & 0 \\ -2 & -3 & 0 & 1+2+3 \end{pmatrix}$$

Since we removed 2 more rows than columns, nullity is at least 2.

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Finding null vectors

$$\begin{pmatrix} 1+2 & 0 & -1 & -2 \\ 0 & 1+2 & -1 & -2 \\ -2 & -3 & 1+2+3 & 0 \\ -2 & -3 & 0 & 1+2+3 \end{pmatrix}$$

Since we removed 2 more rows than columns, nullity is at least 2. Any null vector (a, b, c) of 1×3 matrix gives null vector (a, b, c, c) of 4×4 matrix. (Remember 1 + 2 = 0.)

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Finding null vectors

$$\begin{pmatrix} 1+2 & 0 & -1 & -2 \\ 0 & 1+2 & -1 & -2 \\ -2 & -3 & 1+2+3 & 0 \\ -2 & -3 & 0 & 1+2+3 \end{pmatrix}$$

Since we removed 2 more rows than columns, nullity is at least 2. Any null vector (a, b, c) of 1×3 matrix gives null vector (a, b, c, c) of 4×4 matrix. (Remember 1 + 2 = 0.) We now have factors $12(1 + 2)^2$. To get the blue factors, now pick 1 as the vertex to be removed!

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Higher dimensions: Codimension-1 spanning tree (Adin)

We will use the weighted simplicial matrix-tree theorem. So first we have to find a codimension-1 spanning tree. But it will be a different tree for each color. For each color's factors, treat that color as "last".

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r = 3 (1-dimensional spanning tree): Start with 1, and attach to every other vertex, except blue vertices. Then use 1 to connect the remaining blue vertices.

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Higher dimensions: Codimension-1 spanning tree (Adin)

We will use the weighted simplicial matrix-tree theorem. So first we have to find a codimension-1 spanning tree. But it will be a different tree for each color. For each color's factors, treat that color as "last".

r = 3 (1-dimensional spanning tree): Start with 1, and attach to every other vertex, except blue vertices. Then use 1 to connect the remaining blue vertices.

r = 4 (2-dimensional spanning tree): Start with 1, and attach to every edge with no blue vertices. Then use 1, and attach to all edges using a blue non-1 vertex with a non-red vertex. Finally use 1 with edges with a blue non-1 vertex with a red non-1 vertex.

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The rest of the proof is similar to our $K_{3,2}$ computation:

Reduce by the spanning tree

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The rest of the proof is similar to our $K_{3,2}$ computation:

- Reduce by the spanning tree
- Factor out individual variables from the rows

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 - null vectors of resulting matrix can be expanded to null vectors of full reduced matrix.